

Primes

A prime number p is a positive integer that has exactly 2 divisors, 1 and p .

Facts about primes:

- Every $n \in \mathbb{N}$ is the product of primes.
- Prime factorization is UNIQUE. (proof: read chap. 7)
[Fundamental Theorem of arithmetic]
- $p \mid ab \Rightarrow p \mid a \vee p \mid b$
- $p \mid b \wedge p \nmid a \Rightarrow p \mid \frac{b}{a} \quad \left(\frac{b}{a} = k \in \mathbb{N} \right)$

$$\bullet p|ab \Rightarrow p|a \vee p|b$$

Proof: $p|ab \Rightarrow ab = mp$

If we factor a and b into primes, p must show up by uniqueness of prime factorization.

$$\Rightarrow p|a \vee p|b$$

$$\bullet p|b \wedge p \nmid a \Rightarrow p \mid \frac{b}{a} \quad \left(\frac{b}{a} = k \in \mathbb{N}\right)$$

proof: $\frac{b}{a} = k \Rightarrow p|ak \Rightarrow \underbrace{p|a}_{\text{false}} \vee p|k \Rightarrow p|k.$

• some other properties can be found in chp. 7.

Fermat Theorem

$$p \text{ prime} \wedge p \nmid a \implies a^{p-1} \equiv 1 \pmod{p}$$

- $p \nmid a \implies \gcd(a, p) = 1$

- Consider set $\{1, 2, 3, \dots, p-1\}$

$\times a$
 \pmod{p}
... permute ...

[idea: $\gcd(a, p) = 1$]

(because $\gcd(a, p) = 1$)

So, $a \cdot (2a) \cdot (3a) \cdot (4a) \dots \cdot [(p-1)a] \equiv 1 \cdot 2 \cdot 3 \dots (p-1) = (p-1)!$

$$a^{p-1} \cdot (p-1)! \equiv (p-1)! \pmod{p}$$

- Now $(p-1)!$ and p are co-prime: $D_p = \{1, p\}$

and $p \nmid (p-1)!$ because p can't divide any $x \in \{1, \dots, p-1\}$

So $(p-1)!$ has an inverse mod $p \implies a^{p-1} \equiv 1 \pmod{p}$

Strengthen:

$$p \text{ prime} \iff \forall a \in \{1, \dots, p-1\}, a^{p-1} \equiv 1 \pmod{p}$$

Idea: To check if a number n is prime, make sure $a^{n-1} \equiv 1 \pmod{n}$ for all $a < n$.

Not better than checking $\{1, \dots, n-1\}$ for divisors!

But it turns out, it has good random behavior:

repeat 100 times

- pick random $a < n$

- if $a^{n-1} \not\equiv 1 \pmod{n}$

return false (n is composite)

return true.

Problem: n might be composite and we still return true because we did not pick the "good" a : $a^{n-1} \not\equiv 1 \pmod{n}$

For most composites, the probability of picking a "bad" a is $\leq \frac{1}{2}$ (see chp. 7). Therefore, the prob. of making wrong decision $\leq \left(\frac{1}{2}\right)^{100}$

Other Problems

- a^{n-1} requires $(n-1)$ multiplication
- a^{n-1} is HUGE!

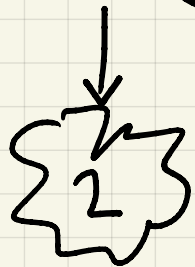
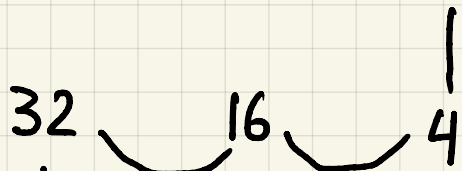
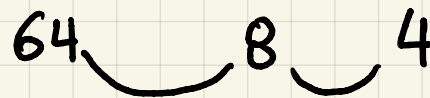
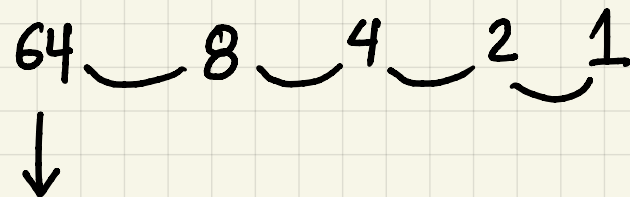
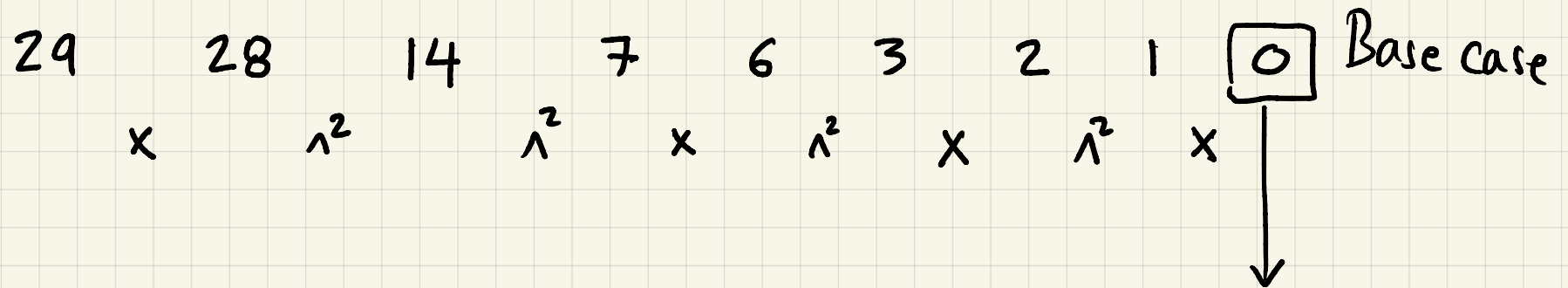
Repeated Squaring:

$$a^b = \begin{cases} 1 & b=0 \\ a \cdot a^{b-1} & b \text{ odd} \\ [a^{b/2}]^2 & b \text{ even} \quad [\text{save mult.}] \end{cases}$$

Combine this with computing everything modulo n on the fly.

Example: $a=2$, $n=30$

Need to find $a^{n-1} = 2^{29}$



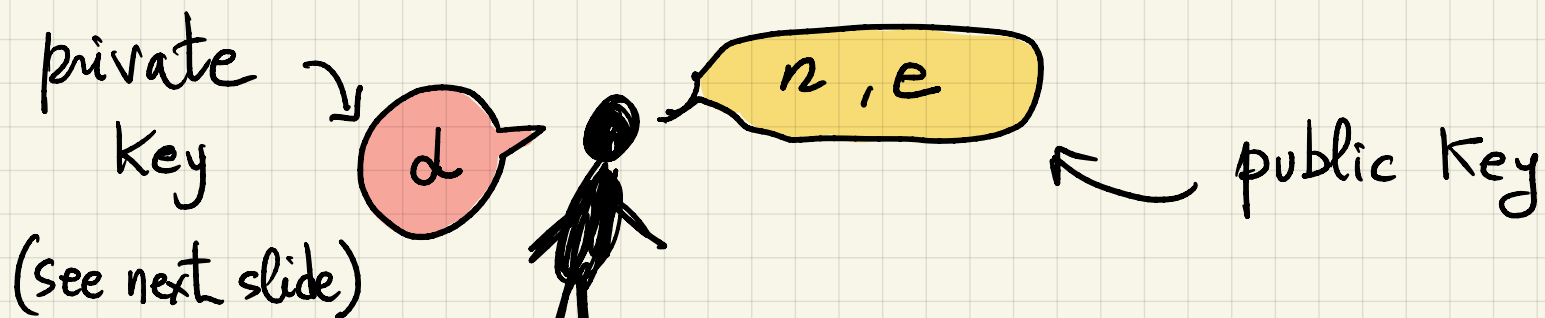
$$\# \text{ mult} \approx 2 \cdot \log_2 b$$

Cryptography

Assume every message is an integer $x < n$.

To send x to person A, send $x^e \bmod n$

where e and n are advertised by A



~~where~~ $n = p \cdot q$ where p, q are large primes

Fact 1: It's hard to factor n into primes, so it's hard to discover p and q

Fact 2: Given $y = x^e \bmod n$, it's hard to figure out x .

Person A also has : $\gcd(e, (p-1)(q-1)) = 1$

so there exists d such that

$$ed \equiv 1 \pmod{(p-1)(q-1)}$$

d can be easily found by A (how?) but not by others.

claim: $y^d \pmod n = x$

$$y^d \equiv (x^e)^d \equiv x^{ed} \equiv x^{(p-1)(q-1)+1} \equiv x \cdot (x^{q-1})^{p-1}$$

• $p \mid x \Rightarrow y^d \equiv x \equiv 0 \pmod p$

• $p \nmid x \Rightarrow p \nmid x^{q-1} \Rightarrow (x^{q-1})^{p-1} \equiv 1 \pmod p$ [Fermat]

$$\Rightarrow y^d \equiv x \pmod p$$

$$y^d \equiv x \pmod{p} \Rightarrow p \mid y^d - x$$

$$y^d \equiv x \pmod{q} \Rightarrow q \mid y^d - x$$

$$n = pq \mid y^d - x \quad (p, q \text{ both prime factors})$$

Therefore $y^d \equiv x \pmod{pq}$

$$y^d \equiv x \pmod{n} \quad \text{😊}$$